Recursively differentiable quasigroups and complete recursive codes

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Abstract. Criteria of recursive differentiability of quasigroups are given. Complete recursive codes which attains the Joshibound are constructed using recursively differentiable k-ary quasigroups.

Keywords: k-recursive code, strong orthogonal system of quasigroups, recursively differentiable quasigroups.

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Let q, n be positive integers and Q be a nonempty set of q elements. A code $C \subseteq Q^n$ of length n over the alphabet Q is called an $[n,k]_Q$-*code* if $|C| = q^k$. An $[n,k,d]_Q$-code is a $[n,k]_Q$-code with the minimal Hamming distance d [1].

According to D.D. Joshi's theorem [2], if C is an $[n,k,d]_Q$-code, then $|C| \leq q^{n-d+1}$, where $|Q| = q$.

If an $[n,k,d]_Q$-code C has the cardinal number $|C| = q^{n-d+1}$ then we say that C *attains the Joshibound*. The problem of description of the parameters q, n and d for which there exist $[n,k,d]_Q$-codes, where $|Q| = q$, attaining the Joshibound is open [1].

It is known that using strong orthogonal systems of k-ary quasigroups ($k \geq 2$), in particular, orthogonal systems of latin squares, such codes can be constructed.

For example, if $\{f_1, f_2, \ldots, f_t\}$, $t \geq 2$, is an orthogonal system of binary quasigroups defined on a set Q of q elements, then

$$C = \{(x, y, f_1(x,y), f_2(x,y), \ldots, f_t(x,y)) \mid x, y \in Q\}$$

is an $[t+2, 2, t+1]_Q$-code, so C attains the Joshibound [2].

This article deals with complete k-recursive codes and recursive differentiability of k-ary quasigroups.

A code C of length n over an alphabet Q is called *complete k-recursive*, where $1 \leq k \leq n$, if there exists a mapping $f : Q^k \longrightarrow Q$ such that every code word $u = (u_0, u_1, \ldots, u_{n-1}) \in C$ satisfies the conditions

$$u_{i+k} = f(u_i, u_{i+1}, \ldots, u_{i+k-1}),$$

for every $i = 0, 1, \ldots, n-k$.

A complete k-recursive code $C \subseteq Q^n$ defined by the mapping f is denoted by $C(n, f)$.

In what follows we will use the notation (x_1^k) for $(x_1, \ldots, x_k)$.

It is proved in [1] and it is easy to see that if $C(n, f)$ is a complete k-recursive code over an alphabet Q then

$$C(n,f) = \{(x_1, \ldots, x_k, f^{(0)}(x_1^{k-1}), \ldots, f^{(n-k-1)}(x_1^k)) \mid x_1, \ldots, x_k \in Q\},$$

where the functions $f^{(0)}, f^{(1)}, \ldots, f^{(n-k-1)}$ are called *k-recursive derivatives* of f and are defined as follows:

$$f^{(0)}(x_1^k) = f(x_1^k),$$
$$f^{(1)}(x_1^k) = f(x_2^k, f^{(0)}(x_1^k)),$$
$$\cdots\cdots\cdots\cdots\cdots\cdots\cdots\cdots\cdots\cdots\cdots\cdots$$
$$f^{(t)}(x_1^k) = f(x_{t+1}^k, f^{(0)}(x_1^k), f^{(1)}(x_1^k), \ldots, f^{(t-1)}(x_1^k)), \quad \text{for} \quad t < k,$$
$$f^{(t)}(x_1^k) = f(f^{(t-k)}(x_1^k), \ldots, f^{(t-1)}(x_1^k)), \quad \text{for} \quad t \geq k.$$

A k-ary quasigroup operation f $(k \geq 2)$ is called *recursively s-differentiable* if its k-recursive derivatives $f^{(0)}, f^{(1)}, \ldots, f^{(s)}$ are k-ary quasigroup operations. Let $k \in \mathbb{N}$, $k \geq 2$, and let $f_1, f_2, \ldots, f_k$ be k-ary operations defined on a set Q. The operations $f_1, f_2, \ldots, f_k$ are called *orthogonal* if the system of equations $\{f_i(x_1, x_2, \ldots, x_k) = a_i\}_{i=1}^k$ has a unique solution for every $a_1, \ldots, a_k \in Q$. It is known and it is easy to see that the k-ary operations $f_1, f_2, \ldots, f_k$, defined on a set Q are orthogonal if and only if the mapping

$$\theta : Q^k \to Q^k, \quad \theta(x_1^k) = (f_1(x_1^k), f_2(x_1^k), \ldots, f_k(x_1^k)) = (f_1, f_2, \ldots, f_k)(x_1^k)$$

is a bijection. In this case we will denote $\theta = (f_1, f_2, \ldots, f_k)$.

A system $\Sigma = \{f_1, f_2, \ldots, f_t\}_{t \geq k}$ of k-ary operations defined on a set Q is called *orthogonal* if every k operations from Σ are orthogonal. A system $\{f_1, f_2, \ldots, f_s\}_{s \geq 1}$ of k-ary operations defined on a set Q is called *strong orthogonal* if the system $\{E_1, \ldots, E_k, f_1, f_2, \ldots, f_s\}$ is orthogonal, where $E_i(x_1^k) = x_i$, for every $(x_1, \ldots, x_k) \in Q^k$ and for every $i = 1, 2, \ldots, k$ (the k-ary selectors).

It follows from the definition that each operation of a strong orthogonal system, which is not a selector, is a quasigroup operation. Every orthogonal system of binary quasigroups is strong orthogonal.

It is proved in [1] that a complete k-recursive code $C(n, f)$ attains the Joshibound if and only if the system of k-recursive derivatives $\{f^{(0)}, f^{(1)}, \ldots, f^{(n-k-1)}\}$ is strong orthogonal. In this case the k-recursive derivatives $f^{(0)}, f^{(1)}, \ldots, f^{(n-k-1)}$ of f are k-ary quasigroup operations, so f is recursively $(n-k-1)$-differentiable. The converse is not true for $k \geq 3$. But for $k = 2$ the following criterion holds.

Proposition 1 ([1]). *A complete 2-recursive code*

$$C(n,f)=\{(x,y,f^{(0)}(x,y),f^{(1)}(x,y),\dots,f^{(n-3)}(x,y))\mid x,y\in Q\}$$

attains the Joshibound if an only if the 2-recursive derivatives $f^{(0)}, f^{(1)}, \dots, f^{(n-3)}$ of f are quasigroup operations.

So a complete 2-recursive code $C(n,f)$ attains the Joshibound if and only if the binary operation f is recursively $(n-3)$-differentiable.

As was announced by G. Belyavskaya in [7] if $Q(f)$ is a binary quasigroup then $f^{(i)}=f\theta^i$, $\forall\, i\in\mathbb{N}$, where θ is the following mapping:

$$\theta: Q^2\longrightarrow Q^2,\ \theta(x,y)=(y,f(x,y)),\quad \forall\,(x,y)\in Q^2.$$

So Proposition 1 has the following algebraic meaning: a binary quasigroup $Q(f)$ is recursively s-differentiable ($s\in\mathbb{N}$) if and only if $f, f\theta,\dots,f\theta^s$, where $\theta=(E_2,f)$, are quasigroup operations. The result announced in [7] is generalized in the following proposition.

Proposition 2. *If f is a k-ary operation ($k\geq 2$) then $f^{(n)}=f\theta^n$ for all $n\in\mathbb{N}$, where*

$$\theta: Q^k\longrightarrow Q^k,\ \theta(x_1^k)=(x_2,\dots,x_k,f(x_1^k)) \tag{1}$$

for every $(x_1^k)\in Q^k$.

Proof: To prove this proposition we will use the mathematical induction.

For $n=0$ and $n=1$, according to the definition of k-recursive derivatives, we have $f^{(0)}=f=f\theta^0$ and $f^{(1)}=f(E_2,\dots,E_k,f)=f\theta$.

Let us suppose that Proposition 2 is true for every n, satisfying the inequalities: $0\leq n\leq s-1<k$. Then for $n=s$, using this assumption, we get:

$$\begin{aligned}f^{(s)}&=f(E_{s+1},\dots,E_k,f^{(0)},\dots,f^{(s-1)})=f(E_{s+1},\dots,E_k,f,f\theta,\dots,f\theta^{s-1})\\&=f(E_s,\dots,E_k,f,f\theta,\dots,f\theta^{s-2})\theta=f\theta^{s-1}\theta=f\theta^s.\end{aligned}$$

For $n=k$ have

$$f^{(k)}=f(f^{(0)},f^{(1)},\dots,f^{(k-1)})=f(E_k,f^{(0)},f^{(1)},\dots,f^{(k-2)})\theta=f\theta^{k-1}\theta=f\theta^k.$$

Let us suppose now that Proposition 2 is true for every $n\leq m-1$, where $m\geq k+1$. Then

$$\begin{aligned}f^{(m)}&=f(f^{(m-k)},\dots,f^{(m-2)},f^{(m-1)})\\&=f(f^{(m-k-1)},\dots,f^{(m-3)},f^{(m-2)})(E_2,\dots,E_k,f)=f\theta^{m-1}\theta=f\theta^m.\end{aligned}$$

So Proposition 2 is true for every $n\in\mathbb{N}$. □

Corollary. *Let $Q(f)$ be an k-ary quasigroup, $k \geq 2$ and $s \in \mathbb{N}$. If $\{f, f\theta, \ldots, f\theta^s\}$, where θ is the mapping defined in (1), is a strong orthogonal system of k-ary operations then $Q(f)$ is recursively s-differentiable.*

As was shown above for $k = 2$ the converse of this corollary is true as well.

Proposition 3. *Let $Q(f)$ be an k-ary quasigroup, $k \geq 2$. Every $k+1$ consecutive k-recursive derivatives $\{f^{(i)}, f^{(i+1)}, \ldots, f^{(i+k)}\}$ of f are orthogonal.*

Proof: If $Q(f)$ is an k-ary quasigroup, $k \geq 2$, then the system $\Sigma = \{E_1, \ldots, E_k, f\}$ is orthogonal, so its subsystem $\{E_2, \ldots, E_k, f\}$ is orthogonal as well, i.e. the mapping

$$\theta : Q^k \longrightarrow Q^k,\ \theta(x_1^k) = (x_2, \ldots, x_k, f(x_1^k)),\quad \forall\, (x_1^k) \in Q^k,$$

is a bijection. Hence each of the following systems is orthogonal:

$$\begin{aligned} \Sigma\theta &= \{E_2, \ldots, E_k, f, f\theta\} = \{E_2, \ldots, E_k, f^{(0)}, f^{(1)}\}, \\ \Sigma\theta^2 &= \{E_3, \ldots, E_k, f, f\theta, f\theta^2\} = \{E_3, \ldots, E_k, f^{(0)}, f^{(1)}, f^{(2)}\}, \ldots, \\ \Sigma\theta^{k-1} &= \{E_k, f, f\theta, \ldots, f\theta^{k-1}\} = \{E_k, f^{(0)}, f^{(1)}, \ldots, f^{(k-1)}\}, \\ \Sigma\theta^k &= \{f, f\theta, \ldots, f\theta^k\} = \{f^{(0)}, f^{(1)}, \ldots, f^{(k)}\} \end{aligned}$$

and

$$\Sigma\theta^s = \{f\theta^{s-k}, \ldots, f\theta^s\} = \{f^{(s-k)}, \ldots, f^{(s)}\},$$

for every $s \geq k + 1$. □

Corollary 1. *A binary quasigroup $Q(f)$ is recursively 1-differentiable if and only if the pair of operations $\{E_1, f^{(1)}\}$ is orthogonal.*

Proof: As $\{E_1, E_2, f\}$ is an orthogonal system, the mapping $\theta = (E_2, f)$ is a bijection and the system $\{E_2, f, f^{(1)}\} = \{E_1, E_2, f\}\theta$ is orthogonal too. Hence, $f^{(1)}$ is a quasigroup operation if and only if the pair $\{E_1, f^{(1)}\}$ is orthogonal. □

Corollary 2. *A ternary quasigroup $Q(f)$ is recursively 1-differentiable iff the systems of ternary operations $\{E_1, E_2, f^{(1)}\}$ and $\{E_1, E_3, f^{(1)}\}$ are orthogonal.*

Let $Q(\cdot)$ be a binary group and let denote by $(\overset{n}{\triangle})$ the n-th 2-recursive derivative of $(\cdot)$, for every $n \in \mathbb{N}$.

Lemma 1. *If $Q(\cdot)$ is an abelian group, then for all $x, y \in Q$ and $n \in \mathbb{N}$ the following equality holds:*

$$x\overset{n}{\triangle}y = x^{b_n}y^{b_{n+1}} \tag{2}$$

where $(b_n)_{n\in\mathbb{N}}$ is the Fibonacci sequence.

Proof: We will use the mathematical induction.

For $n = 0$ have $x\overset{0}{\triangle}y = x \cdot y$ so $x\overset{0}{\triangle}y = x^{b_0} \cdot y^{b_1}$.

For $n = 1$ have $x\overset{1}{\triangle}y = y \cdot xy = x \cdot y^2 = x^{b_1} \cdot y^{b_2}$.

Suppose that Lemma 1 is true for every $n \le k$. Using this assumption and the definition of the Fibonacci sequence, for $n = k+1$ we get

$$\begin{aligned} x \overset{k+1}{\triangle} y = (x \overset{k-1}{\triangle} y)(x\overset{k}{\triangle}y) &= x^{b_{k-1}} \cdot y^{b_k} \cdot x^{b_k} \cdot y^{b_{k+1}} \\ &= x^{b_{k-1}+b_k} \cdot y^{b_k+b_{k+1}} = x^{b_{k+1}} \cdot y^{b_{k+2}}. \end{aligned}$$

So the equality (2) is true for every $x, y \in Q$ and for every $n \in \mathbb{N}$. □

Theorem 1. *A binary abelian group $Q(\cdot)$ is recursively s-differentiable, where $s \ge 1$, if and only if the mappings $x \mapsto x^{b_i}$, where $(b_n)_{n\in\mathbb{N}}$ is the Fibonacci sequence, are bijections for all $i \in \{0, 1, 2, \dots, s+1\}$.*

Proof: According to the definition a group $Q(\cdot)$ is recursively s-differentiable if and only if its 2-recursive derivatives $(\overset{1}{\triangle}), (\overset{2}{\triangle}), \dots, (\overset{s}{\triangle})$ are quasigroup operations. Hence $Q(\cdot)$ is recursively s-differentiable if and only if each of the equations $x\overset{i}{\triangle}a = c$, $a\overset{i}{\triangle}y = c$, $i \in \{0, 1, 2, \dots, s\}$, has a unique solution for every $a, c \in Q$. Now, using the equalities (2) we get: $x\overset{i}{\triangle}a = c \Leftrightarrow x^{b_i} \cdot a^{b_{i+1}} = c \Leftrightarrow x^{b_i} = c \cdot a^{-b_{i+1}}$ and $a\overset{i}{\triangle}y = c \Leftrightarrow y^{b_{i+1}} \cdot a^{b_i} = c \Leftrightarrow y^{b_{i+1}} = c \cdot a^{-b_i}$ for every $a, c \in Q$ and for every $i \in \{0, 1, 2, \dots, s\}$. So $(\overset{1}{\triangle}), (\overset{2}{\triangle}), \dots, (\overset{s}{\triangle})$ are quasigroup operations if and only if the mappings $x \mapsto x^{b_i}$, where $(b_i)_{i\in\mathbb{N}}$ is the Fibonacci sequence, are bijections for every $i \in \{0, 1, 2, \dots, s+1\}$. □

Proposition 4. *If $Q(\cdot)$ is an arbitrary recursively s-differentiable binary group, where $s \ge 1$, then the mappings $x \mapsto x^{b_i}$, where $(b_i)_{i\in\mathbb{N}}$ is the Fibonacci sequence, are bijections for all $i \in \{0, 1, 2, \dots, s+1\}$.*

Proof: If $Q(\cdot)$ is recursively s-differentiable, with unit e, then each of the equations $e\overset{i}{\triangle}x = c$ and $y\overset{i}{\triangle}e = c$, $i \in \{0, 1, 2, \dots, s\}$, has a unique solution. So as

$e\overset{i}{\triangle}x = c \Leftrightarrow x^{b_{i+1}} = c$ and $y\overset{i}{\triangle}e = c \Leftrightarrow y^{b_i} = c$, we get that each of the mappings $x \mapsto x^{b_i}$, $i \in \{0, 1, 2, \ldots, s+1\}$, is a bijection. □

When $s = 1$ Theorem 1 is true for an arbitrary binary group as we can see from the following proposition.

Proposition 5. *A binary group $Q(\cdot)$ is recursively 1-differentiable if and only if the mapping $z \mapsto z^2$ is a bijection.*

PROOF: According to the definition, a binary group $Q(\cdot)$ is recursively 1-differentiable if and only if its 2-recursive derivative $(\overset{1}{\triangle})$ is a quasigroup operation. So as $a\overset{1}{\triangle}x = b \Leftrightarrow x \cdot ax = b \Leftrightarrow xaxa = ba \Leftrightarrow (xa)^2 = ba$, for every $a, b \in Q$, we get that the mapping $z \mapsto z^2$ is a bijection if and only if the equation $a\overset{1}{\triangle}(za^{-1}) = b$ has a unique solution z for every $a, b \in Q$.

From the equivalences $x\overset{1}{\triangle}a = b \Leftrightarrow a \cdot xa = b \Leftrightarrow x = a^{-1}ba^{-1}$ it follows that in a binary quasigroup $Q(\cdot)$ the equation $x\overset{1}{\triangle}a = b$ has always a unique solution for every $a, b \in Q$. So if $Q(\cdot)$ is a group then $Q(\overset{1}{\triangle})$ is a quasigroup if and only if the mapping $z \mapsto z^2$ is a bijection. □

Corollary. *A finite binary group is recursively 1-differentiable if and only if it is of odd order.*

Indeed, it is known [3] that a finite group is of odd order if and only if the mapping $z \mapsto z^2$ is a bijection.

Proposition 6. *If $Q(\cdot)$ is a binary group with unit e, then $Q(\overset{1}{\triangle})$ is a semigroup if and only if $x^2 = e$, for every $x \in Q$.*

PROOF: So as $(x\overset{1}{\triangle}y)\overset{1}{\triangle}z = zyxyz$ and $x\overset{1}{\triangle}(y\overset{1}{\triangle}z) = zyzxzyz$, for all $x, y, z \in Q$, we get that the operation $(\overset{1}{\triangle})$ is associative if and only if $x = zxz$, for every $x, z \in Q$. Taking $x = e$ in the last equality we get $z^2 = e$, for all $z \in Q$. Conversely, if $x^2 = e$, for all $x \in Q$, then $x = x^{-1}$ and $xz \cdot xz = e$, $\forall\, x, z \in Q$, so $zxz = x^{-1} = x$, for all $x, z \in Q$, i.e. $(\overset{1}{\triangle})$ is associative. □

Corollary. *If $Q(\cdot)$ is a nontrivial recursively 1-differentiable group then its 2-recursive derivative $Q(\overset{1}{\triangle})$ cannot be a group.*

PROOF: Indeed, if $Q(\cdot)$ is recursively 1-differentiable and $Q(\overset{1}{\triangle})$ is a group, then according to Proposition 5, the mapping $z \mapsto z^2$ is a bijection and by Proposition 6 we get $|Q| = 1$. □

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